

**CS 441 Discrete Mathematics for CS**  
**Lecture 3**

**Propositional logic: Equivalences**  
**Predicate logic**

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**Propositional logic: review**

- **Propositional logic:** a formal language for making logical inferences
- A **proposition** is a statement that is either true or false.
- A **compound proposition** can be created from other propositions using logical connectives
- **The truth of a compound proposition** is defined by truth values of elementary propositions and the meaning of connectives.
- **The truth table for a compound proposition:** table with entries (rows) for all possible combinations of truth values of elementary propositions.

## Tautology and Contradiction

### Definitions:

- A compound proposition that is always true for all possible truth values of the propositions is called a **tautology**.
- A compound proposition that is always false is called a **contradiction**.
- A proposition that is neither a tautology nor contradiction is called a **contingency**.

Example:  $p \vee \neg p$  is a **tautology**.

$p$	$\neg p$	$p \vee \neg p$
T	F	T
F	T	T

## Tautology and Contradiction

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Example:  $p \wedge \neg p$  is a **contradiction**.

$p$	$\neg p$	$p \wedge \neg p$
T	F	F
F	T	F

## Equivalence

- Some propositions may be equivalent. Their truth values in the truth table are the same.
- Example:  $p \rightarrow q$  is equivalent to  $\neg q \rightarrow \neg p$  (**contrapositive**)

$p$	$q$	$p \rightarrow q$	$\neg q \rightarrow \neg p$
T	T		
T	F		
F	T		
F	F		

## Equivalence

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$p$	$q$	$p \rightarrow q$	$\neg q \rightarrow \neg p$
T	T	T	T
T	F	F	F
F	T	T	T
F	F	T	T

- **Equivalent statements** are important for logical reasoning since they can be substituted and can help us to make a logical argument.

## Logical equivalence

- **Definition:** The propositions  $p$  and  $q$  are called **logically equivalent** if  $p \leftrightarrow q$  is a tautology (alternately, if they have the same truth table). The notation  $p \Leftrightarrow q$  denotes  $p$  and  $q$  are logically equivalent.

### Examples of equivalences:

- **DeMorgan's Laws:**

- 1)  $\neg(p \vee q) \Leftrightarrow \neg p \wedge \neg q$
- 2)  $\neg(p \wedge q) \Leftrightarrow \neg p \vee \neg q$

## Equivalence

### Example of important equivalences

- **DeMorgan's Laws:**

- 1)  $\neg(p \vee q) \Leftrightarrow \neg p \wedge \neg q$
- 2)  $\neg(p \wedge q) \Leftrightarrow \neg p \vee \neg q$

Use the truth table to prove that the two propositions are logically equivalent

$p$	$q$	$\neg p$	$\neg q$	$\neg(p \vee q)$	$\neg p \wedge \neg q$
T	T	F	F	<b>F</b>	<b>F</b>
T	F	F	T	<b>F</b>	<b>F</b>
F	T	T	F	<b>F</b>	<b>F</b>
F	F	T	T	<b>T</b>	<b>T</b>

## Important logical equivalences

- **Identity**

- $p \wedge T \Leftrightarrow p$

- $p \vee F \Leftrightarrow p$

- **Domination**

- $p \vee T \Leftrightarrow T$

- $p \wedge F \Leftrightarrow F$

- **Idempotent**

- $p \vee p \Leftrightarrow p$

- $p \wedge p \Leftrightarrow p$

## Important logical equivalences

- **Double negation**

- $\neg(\neg p) \Leftrightarrow p$

- **Commutative**

- $p \vee q \Leftrightarrow q \vee p$

- $p \wedge q \Leftrightarrow q \wedge p$

- **Associative**

- $(p \vee q) \vee r \Leftrightarrow p \vee (q \vee r)$

- $(p \wedge q) \wedge r \Leftrightarrow p \wedge (q \wedge r)$

## Important logical equivalences

- **Distributive**

- $p \vee (q \wedge r) \Leftrightarrow (p \vee q) \wedge (p \vee r)$
- $p \wedge (q \vee r) \Leftrightarrow (p \wedge q) \vee (p \wedge r)$

- **De Morgan**

- $\neg(p \vee q) \Leftrightarrow \neg p \wedge \neg q$
- $\neg(p \wedge q) \Leftrightarrow \neg p \vee \neg q$

- **Other useful equivalences**

- $p \vee \neg p \Leftrightarrow T$
- $p \wedge \neg p \Leftrightarrow F$
- $p \rightarrow q \Leftrightarrow (\neg p \vee q)$

## Using logical equivalences

- **Equivalences can be used in proofs.** A proposition or its part can be transformed using equivalences and some conclusion can be reached.

**Example:** Show that  $(p \wedge q) \rightarrow p$  is a tautology.

- Proof: (we must show  $(p \wedge q) \rightarrow p \Leftrightarrow T$ )

$$\begin{aligned}(p \wedge q) \rightarrow p &\Leftrightarrow \neg(p \wedge q) \vee p && \text{Useful} \\ &\Leftrightarrow [\neg p \vee \neg q] \vee p && \text{DeMorgan} \\ &\Leftrightarrow [\neg q \vee \neg p] \vee p && \text{Commutative} \\ &\Leftrightarrow \neg q \vee [\neg p \vee p] && \text{Associative} \\ &\Leftrightarrow \neg q \vee [T] && \text{Useful} \\ &\Leftrightarrow T && \text{Domination}\end{aligned}$$

## Using logical equivalences

- **Equivalences can be used in proofs.** A proposition or its part can be transformed using equivalences and some conclusion can be reached.

**Example:** Show  $(p \wedge q) \rightarrow p$  is a tautology.

- Alternative proof:

p	q	$p \wedge q$	$(p \wedge q) \rightarrow p$
T	T	T	T
T	F	F	T
F	T	F	T
F	F	F	T

## Using logical equivalences

- **Equivalences can be used in proofs.** A proposition or its part can be transformed using equivalences and some conclusion can be reached.
- **Proofs that rely on logical equivalences can replace truth table approach**
  - **Why?**
  - The truth table has  $2^n$  rows, where  $n$  is the number of elementary propositions
  - If  $n$  is large building the truth table may become infeasible

## Using logical equivalences

- Equivalences can be used in proofs. A proposition or its part can be transformed using equivalences and some conclusion can be reached.
- **Example 2:** Show  $(p \rightarrow q) \Leftrightarrow (\neg q \rightarrow \neg p)$

**Proof:**

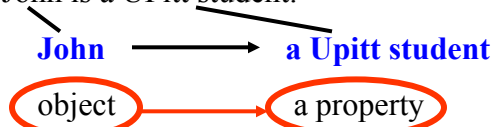
- $(p \rightarrow q) \Leftrightarrow (\neg q \rightarrow \neg p)$
- $\Leftrightarrow \neg(\neg q) \vee (\neg p)$  Useful
- $\Leftrightarrow q \vee (\neg p)$  Double negation
- $\Leftrightarrow \neg p \vee q$  Commutative
- $\Leftrightarrow p \rightarrow q$  Useful

**End of proof**

## Limitations of the propositional logic

- **Propositional logic:** the world is described in terms of propositions
- **A proposition** is a statement that is either true or false.
- **Limitations:**
  - **objects** in elementary statements, their properties and relations are not explicitly represented in the propositional logic
- **Example:**

– “John is a UPitt student.”



– Objects and properties are hidden in the statement, it is not possible to reason about them

## Limitations of the propositional logic

- **Statements for groups of objects**
  - In propositional logic these must be exhaustively enumerated
- **Example:**
  - If John is a CS UPitt graduate then John has passed cs441

**Translation:**

  - John is a CS UPitt graduate  $\rightarrow$  John has passed cs441

Similar statements can be written for other Upitt graduates:

  - Ann is a CS Upitt graduate  $\rightarrow$  Ann has passed cs441
  - Ken is a CS Upitt graduate  $\rightarrow$  Ken has passed cs441
  - ...
- **What is a more natural solution to express the above knowledge?**

## Limitations of the propositional logic

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  - In propositional logic these must be exhaustively enumerated
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**Translation:**

  - John is a CS UPitt graduate  $\rightarrow$  John has passed cs441

Similar statements can be written for other Upitt graduates:

  - Ann is a CS Upitt graduate  $\rightarrow$  Ann has passed cs441
  - Ken is a CS Upitt graduate  $\rightarrow$  Ken has passed cs441
  - ...
- **Solution:** make statements with **variables**
  - If  $x$  is a CS Upitt graduate then  $x$  has passed cs441
  - $x$  is a CS UPitt graduate  $\rightarrow$   $x$  has passed cs441

## Predicate logic

### Remedies the limitations of the propositional logic

- Explicitly models objects and their properties
- Allows to make statements with variables and quantify them

### Basic building blocks of the predicate logic:

- **Constant** –models a specific object

**Examples:** “John”, “France”, “7”

- **Variable** – represents object of specific type (**defined by the universe of discourse**)

**Examples:**  $x, y$

(universe of discourse can be people, students, numbers)

- **Predicate** - over one, two or many variables or constants.
  - Represents properties or relations among objects

**Examples:** Red(car23), student( $x$ ), married(John,Ann)

## Predicates

**Predicates** represent properties or relations among objects

A predicate  $P(x)$  assigns a value **true or false** to each  $x$  depending on whether the property holds or not for  $x$ .

- The assignment is best viewed as a big table with the variable  $x$  substituted for objects from the universe of discourse

### Example:

- Assume **Student( $x$ )** where the universe of discourse are people
- Student(John) .... T (if John is a student)
- Student(Ann) .... T (if Ann is a student)
- Student(Jane) ..... F (if Jane is not a student)
- ...

## Predicates

Assume a predicate  $P(x)$  that represents the statement:

- $x$  is a prime number

What are the truth values of:

- $P(2)$  T
- $P(3)$  T
- $P(4)$  F
- $P(5)$  T
- $P(6)$  F
- $P(7)$  T

All statements  $P(2)$ ,  $P(3)$ ,  $P(4)$ ,  $P(5)$ ,  $P(6)$ ,  $P(7)$  are propositions

## Predicates

Assume a predicate  $P(x)$  that represents the statement:

- $x$  is a prime number

What are the truth values of:

- $P(2)$  T
- $P(3)$  T
- $P(4)$  F
- $P(5)$  T
- $P(6)$  F
- $P(7)$  T

Is  $P(x)$  a proposition? No. Many possible substitutions are possible.

## Predicates

### Important:

- predicate  $P(x)$  is **not a proposition** since there are more objects it can be applied to

**This is the same as in propositional logic ...**

### ... But the difference is:

- propositional logic does not let us go inside the statements and manipulate  $x$
- predicate logic allows us to explicitly manipulate and substitute for the objects

## Predicates

- Predicates can have **more arguments** which represent the **relations between objects**

### Example:

- $\text{Older}(\text{John}, \text{Peter})$  denotes ‘John is older than Peter’
  - this is a proposition because it is either true or false
- $\text{Older}(x,y)$  - ‘ $x$  is older than  $y$ ’
  - not a proposition, but after the substitution it becomes one

## Predicates

- Predicates can have **more arguments** which represent the **relations between objects**

### Example:

- Let  $Q(x,y)$  denote ' $x+5 > y$ '
  - Is  $Q(x,y)$  a proposition? **No!**
  - Is  $Q(3,7)$  a proposition? **Yes.** It is true.
  - What is the truth value of
    - $Q(3,7)$  T
    - $Q(1,6)$  F
    - $Q(2,2)$  T
  - Is  $Q(3,y)$  a proposition? **No!** We cannot say if it is true or false.

## Quantified statements

**Predicate logic allows us to make statements about groups of objects**

- To do this we use special quantified expressions

Two types of quantified statements:

- **universal**

**Example:** 'all CS Upitt graduates have to pass cs441''

- the statement is true for all graduates

- **existential**

**Example:** 'Some CS Upitt students graduate with honor.'

- the statement is true for some people

## Universal quantifier

**Defn:** The universal quantification of  $P(x)$  is the proposition " $P(x)$  is true for all values of  $x$  in the universe of discourse." The notation  $\forall x P(x)$  denotes the universal quantification of  $P(x)$ , and is expressed as **for every  $x$ ,  $P(x)$** .

**Example:**

- Let  $P(x)$  denote  $x > x - 1$ .
- What is the truth value of  $\forall x P(x)$ ?
- Assume the universe of discourse of  $x$  is all real numbers.
- **Answer:** Since every number  $x$  is greater than itself minus 1. Therefore,  **$\forall x P(x)$  is true.**

## Universal quantifier

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**Example 2:**

- Let  $T(x)$  denote  $x > 5$ .
- What is the truth value of  $\forall x T(x)$ ?
- Assume the universe of discourse of  $x$  is all real numbers.
- **Answer:**
  - Since  $3 > 5$  is false. So,  $T(x)$  is not true for all values of  $x$ . Therefore, it is **false that  $\forall x T(x)$** .

## Universal quantifier

**Quantification converts** a propositional function into a **proposition** by binding a variable to a set of values from the universe of discourse.

### Example:

- Let  $P(x)$  denote  $x > x - 1$ .
- Is  $P(x)$  a proposition? **No**. Many possible substitutions.
- Is  $\forall x P(x)$  a proposition? **Yes**. True if for all  $x$  from the universe of discourse  $P(x)$  is true. Which holds?
  
- Is  $\forall x Q(x,y)$  a proposition? **No**. The variable  $y$  is free and can be substituted by many objects.