CS 1571 Introduction to AI Lecture 8

Constraint satisfaction search. Combinatorial optimization search.

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Constraint satisfaction problem (CSP)

Objective:

- Find a configuration satisfying goal conditions
- Constraint satisfaction problem (CSP) is a configuration search problem where:
 - A state is defined by a set of variables and their values
 - Goal condition is represented by a set constraints on possible variable values

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CSP example: N-queens

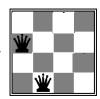
Goal: n queens placed in non-attacking positions on the board

Variables:

• Represent queens, one for each column:

$$-Q_1,Q_2,Q_3,Q_4$$

- Values:
 - Row placement of each queen on the board{1, 2, 3, 4}



$$Q_1 = 2, Q_2 = 4$$

Constraints: $Q_i \neq Q_j$ Two queens not in the same row $|Q_i - Q_j| \neq |i - j|$ Two queens not on the same diagonal

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Map coloring

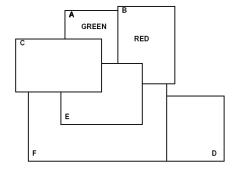
Color a map using k different colors such that no adjacent countries have the same color

Variables:

• Represent countries

$$-A,B,C,D,E$$

- Values:
 - K -different colors{Red, Blue, Green,...}



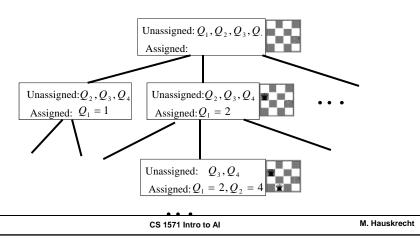
Constraints: $A \neq B, A \neq C, C \neq E$, etc

An example of a problem with binary constraints

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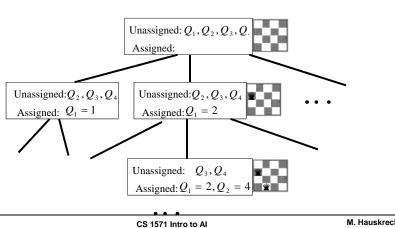
Solving a CSP through standard search

- Maximum depth of the tree: Number of variables in the CSP
- **Depth of the solution:** Number of variables in the CSP
- **Branching factor:** if we fix the order of variable assignments the branch factor depends on the number of their values



Solving a CSP through standard search

- What search algorithm to use: Depth first search !!!
 - Since we know the depth of the solution
 - We do not have to keep large number of nodes in queues



Constraint consistency

Question:

- When to check the constraints defining the goal condition?
- The violation of constraints can be checked:
 - at the end (for the leaf nodes)
 - for each node of the search tree during its generation or before its expansion

Checking the constraints for intermediate nodes:

• More efficient: cuts branches of the search tree early

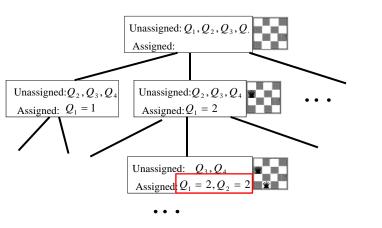
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Constraint consistency

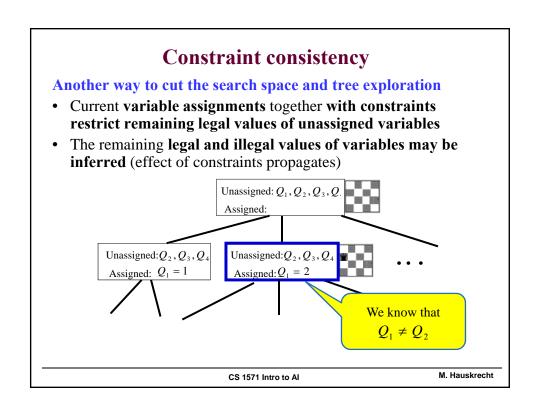
Checking the constraints for intermediate nodes:

• More efficient: cuts branches of the search tree early



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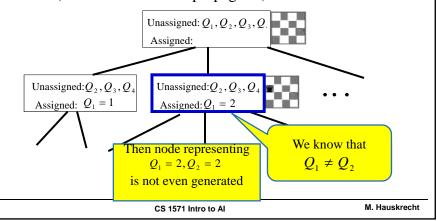
Checking the constraints for intermediate nodes: • More efficient: cuts branches of the search tree early Unassigned: Q_1, Q_2, Q_3, Q_4 Assigned: Q_1, Q_2, Q_3, Q_4



Constraint consistency

Another way to cut the search space and tree exploration

- Current variable assignments together with constraints restrict remaining legal values of unassigned variables
- The remaining legal and illegal values of variables may be inferred (effect of constraints propagates)



Constraint propagation

A state (more broadly) is defined:

- by a set of assigned variables, their values and
- a list of legal and illegal assignments for unassigned variables Legal and illegal assignments can be represented:
- equations (value assignments) and
- disequations (list of invalid assignments)

$$A = \text{Red}$$
, Blue $C \neq \text{Red}$

Constraints + assignments

can entail new equations and disequations

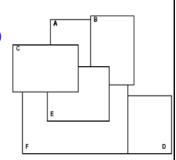
$$A = \text{Red} \rightarrow B \neq \text{Red}$$

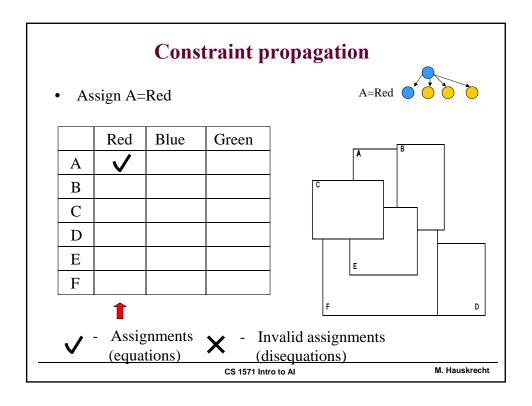
Constraint propagation: the process

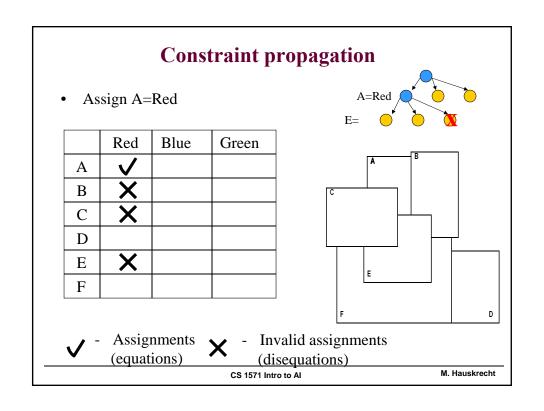
of inferring of new equations and disequations

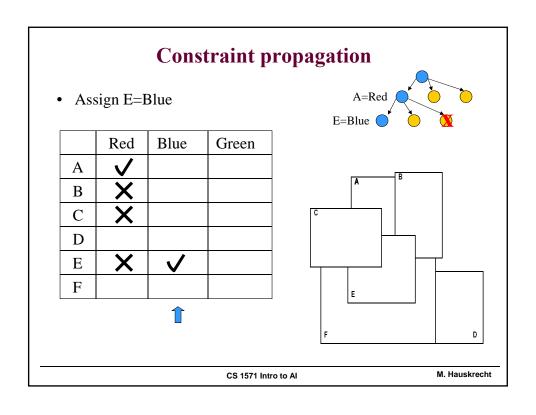
from existing equations and disequations

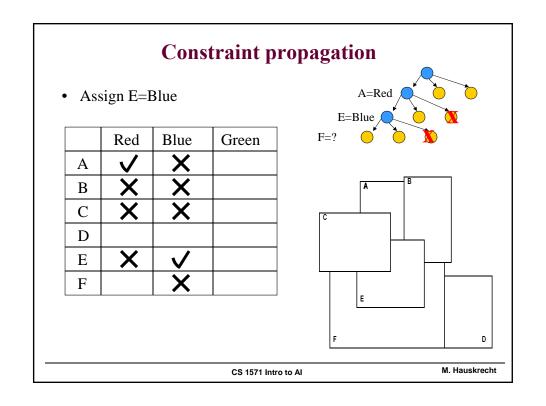
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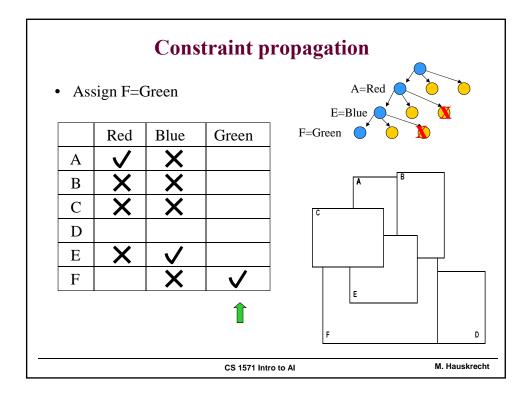












Constraint propagation

Three techniques for propagating the effects of past assignments and constraints:

- Node propagation
- Arc consistency
- Forward checking
- Difference:
 - Completeness of inferences
 - Time complexity of inferences.

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Constraint propagation

- 1. Node consistency. Infers:
 - equations (valid assignments) or disequations (invalid assignments) for an individual variable by applying a unary constraint
- 2. Arc consistency. Infers:
 - disequations from the set of equations and disequations defining the partial assignment, and a constraint
 - equations through the exhaustion of alternatives
- 3. Forward checking. Infers:
 - disequations from a set of equations defining the partial assignment, and a constraint
 - Equations through the exhaustion of alternatives

Restricted forward checking:

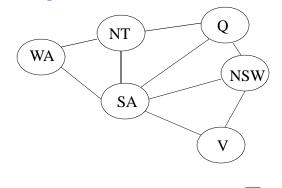
uses only active constraints (active constraint – only one variable unassigned in the constraint)

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Example

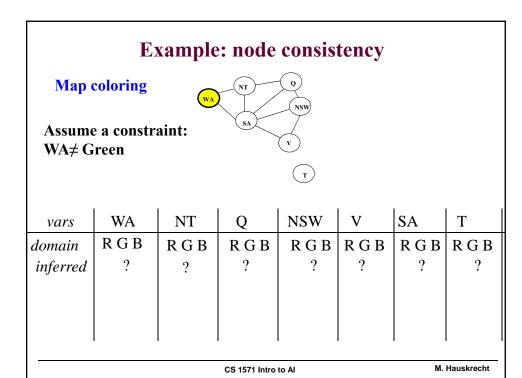
Map coloring of Australia territories

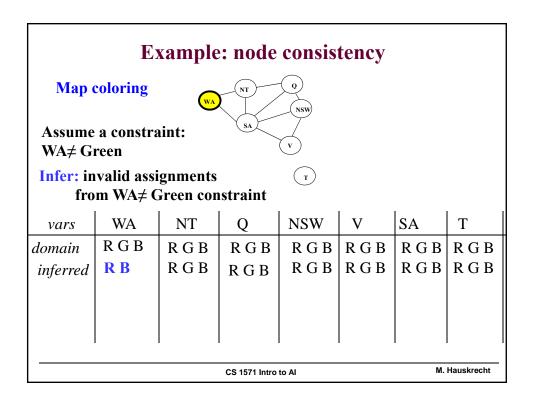


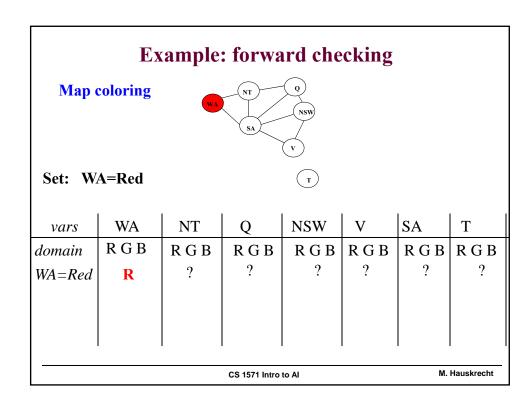
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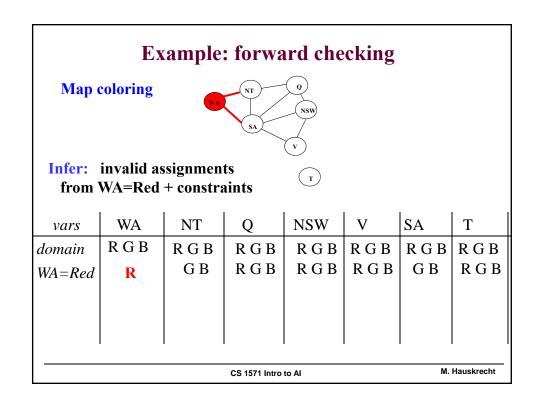
Note: problem with binary constraints

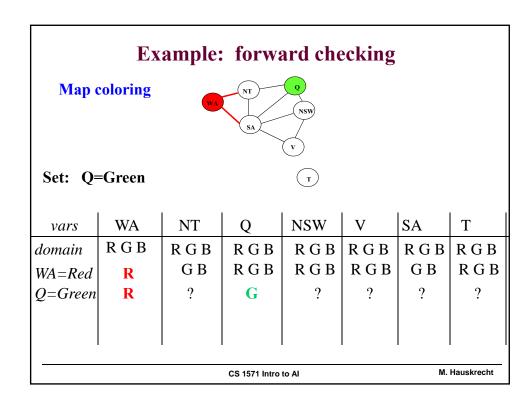
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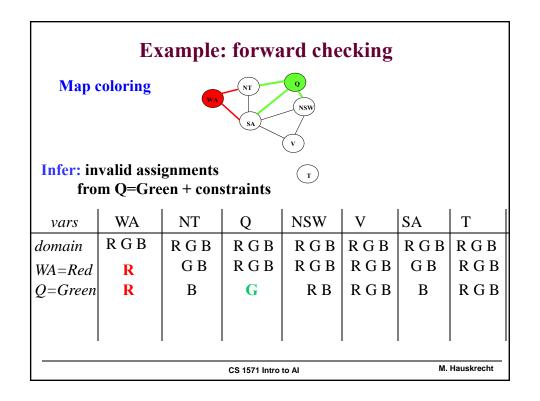






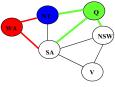








Map coloring



Infer: NT=B

Exhaustions of alternatives

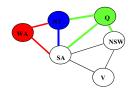
vars	WA	NT	Q	NSW	V	SA	T
domain	R G B	R G B	R G B	R G B	R G B	R G B	R G B
WA=Red	R	G B	R G B	R G B	R G B	G B	R G B
Q=Green	R	В	\mathbf{G}	R B	RGB	В	RGB
Infer NT	R	В	\mathbf{G}	?	?	?	?

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Example: forward checking

Map coloring

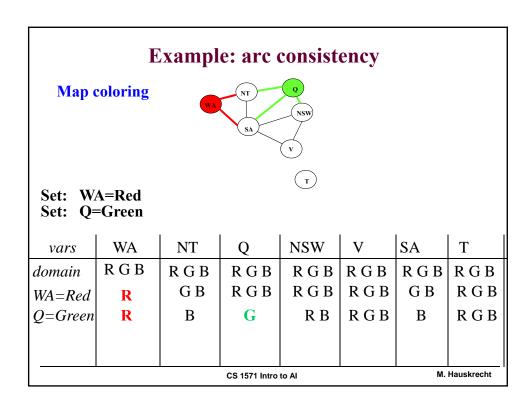


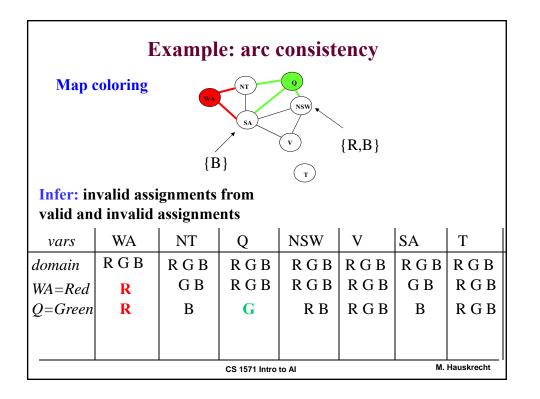
Infer: invalid assignments

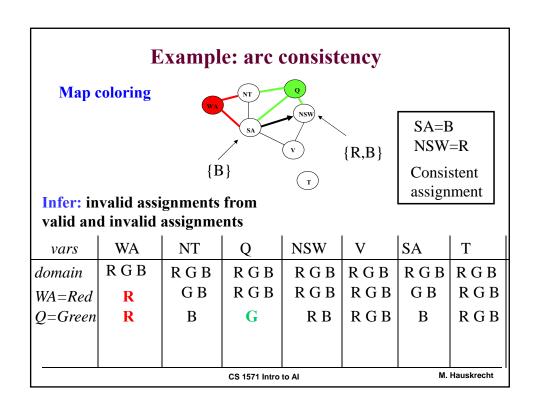
from NT=B

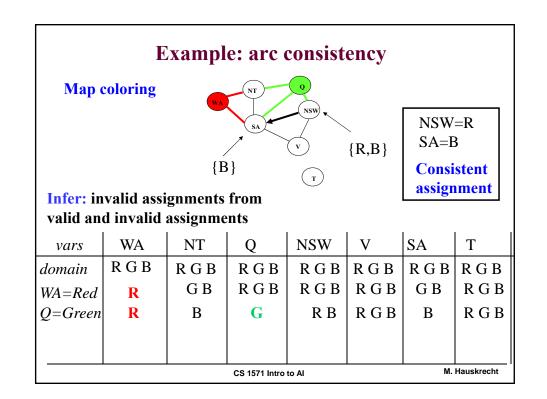
vars	WA	NT	Q	NSW	V	SA	T
domain	R G B	R G B	RGB	R G B	R G B	R G B	R G B
WA=Red	R	G B	RGB	R G B	R G B	G B	R G B
Q=Green	R	В	G	R B	R G B	В	RGB
Infer NT	R	B	G	R B	R G B		RGB

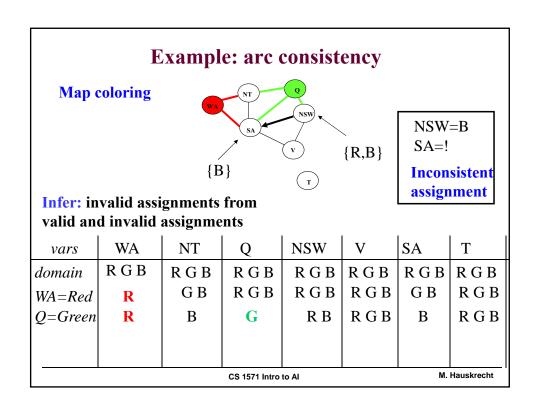
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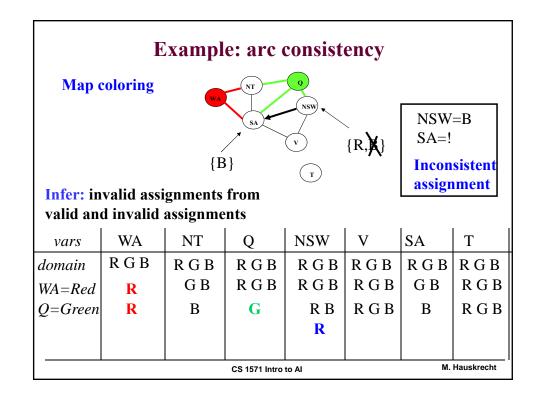












Heuristics for CSPs

CSP searches the space in the depth-first manner.

But we still can choose:

- Which variable to assign next?
- Which value to choose first?

Heuristics

- Most constrained variable
 - Which variable is likely to become a bottleneck?
- Least constraining value
 - Which value gives us more flexibility later?

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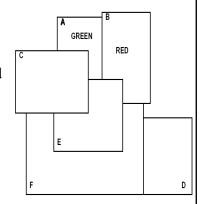
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Heuristics for CSP

Examples: map coloring

Heuristics

- Most constrained variable
 - Country E is the most constrained one (cannot use Red, Green)
- Least constraining value
 - Assume we have chosen variable C
 - Red is the least constraining valid color for the future



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Finding optimal configurations

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Search for the optimal configuration

Constrain satisfaction problem:

Objective: find a configuration that satisfies all constraints



Optimal configuration problem:

Objective: find the best configuration

The quality of a configuration: is defined by some quality measure that reflects our preference towards each configuration (or state)

Our goal: optimize the configuration according to the quality measure also referred to as **objective function**

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Search for the optimal configuration

Optimal configuration search:

- Configurations are described in terms of variables and their values
- Each configuration has a quality measure
- Goal: find the configuration with the best value

If the space of configurations we search among is

- Discrete or finite
 - then it is a combinatorial optimization problem
- Continuous
 - then it is a parametric optimization problem

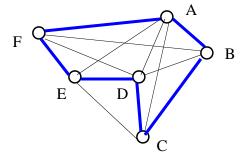
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Example: Traveling salesman problem

Problem:

- A graph with distances
- A tour a path that visits every city once and returns to the start e.g. ABCDEF

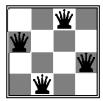


• Goal: find the shortest tour

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Example: N queens

- A CSP problem
- Is it possible to formulate the problem as an optimal configuration search problem ?

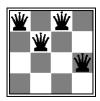


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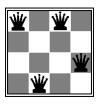
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Example: N queens

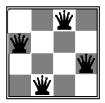
- A CSP problem
- Is it possible to formulate the problem as an optimal configuration search problem? Yes.
- The quality of a configuration in a CSP can be measured by the number of violated constraints
- Solving: minimize the number of constraint violations



of violations =3



of violations =1



of violations =0

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Iterative optimization methods

- Searching systematically for the best configuration with the DFS may not be the best solution
- Worst case running time:
 - Exponential in the number of variables
- Solutions to **large 'optimal' configuration** problems are often found using iterative optimization methods
- Methods:
 - Hill climbing
 - Simulated Annealing
 - Genetic algorithms

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Iterative optimization methods

Properties:

- Search the space of "complete" configurations
- Take advantage of local moves
 - Operators make "local" changes to "complete" configurations
- Keep track of just one state (the current state)
 - no memory of past states
 - !!! No search tree is necessary !!!

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Example: N-queens

- "Local" operators for generating the next state:
 - Select a variable (a queen)
 - Reallocate its position



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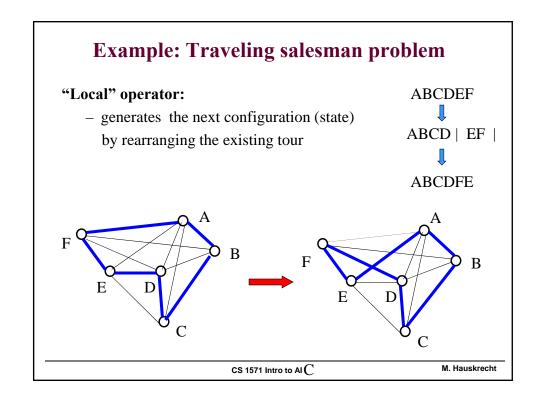
Example: Traveling salesman problem

"Local" operator for generating the next state:

- divide the existing tour into two parts,
- reconnect the two parts in the opposite order

Example: ABCDEF F CS 1571 Intro to AI A M. Hauskrecht

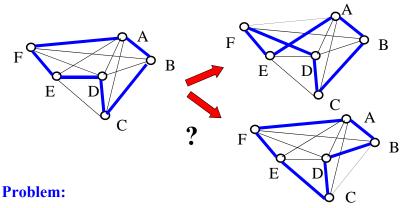
Example: Traveling salesman problem "Local" operator for generating the next state: • divide the existing tour into two parts, • reconnect the two parts in the opposite order Example: ABCDEF ABCDEF ABCDEF ABCDFE ABCDFE CS 1571 Intro to Al M. Hauskrecht



Searching the configuration space

Search algorithms

• keep only one configuration (the current configuration)



• How to decide about which operator to apply?

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Search algorithms

Strategies to choose the configuration (state) to be visited next:

- Hill climbing
- Simulated annealing
- Later: Extensions to multiple current states:
 - Genetic algorithms
- **Note:** Maximization is inverse of the minimization

$$\min f(X) \Leftrightarrow \max \left[-f(X) \right]$$

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