

CS 1510 Midterm 2  
Fall 2007

1. (20 points) Consider the following problem. The input consists of  $n$  positive integers  $V = \{v_1, \dots, v_n\}$ . Let  $L = \sum_{i=1}^n v_i$ . The problem is to determine if there are three disjoint subsets  $S$ ,  $P$  and  $T$  of  $V$  such that

$$\sum_{v_i \in S} v_i = \prod_{v_i \in P} v_i = \sum_{v_i \in T} v_i^3$$

That is, no element of  $V$  can be in more than one of  $S$ ,  $P$  and  $T$ . You want the sum of the numbers in  $S$  to be equal to the product of the numbers in  $P$ , and equal to the sum of the cubes of the numbers in  $T$ . It is okay for a number to be in none of  $S$ ,  $P$  or  $T$ . Give an algorithm for this problem whose running time is polynomial in  $n$  and  $L$ .

2. (20 points)
- (a) Explain how to solve the Parallel Prefix Problem on an EREW machine in time  $O(\log n)$  on input of size  $n$  with  $n$  processors. Start with defining the Parallel Prefix Problem.
  - (b) What is the efficiency of the algorithm in the previous subproblem? Start with a definition of efficiency. You need not have solved the previous subproblem to answer this question.
  - (c) What would the Folding Principle say about the time for the algorithm in the first subproblem if there were only  $n^{1/4}$  processors? Start with a definition of the Folding Principle. You need not have solved the first subproblem to answer this question.
  - (d) Explain how to solve the Parallel Prefix Problem on an EREW machine in time  $O(\log n)$  on input of size  $n$  with  $n/\log n$  processors.

3. (20 points)
- (a) Explain how to find the maximum of  $n$  numbers in time  $O(1)$  on an CRCW-common machine with  $n^2$  processors.
  - (b) Explain how to find the maximum of  $n$  numbers in time  $O(\log \log n)$  on an CRCW-common machine with  $n$  processors.

4. (45 points)
- (a) Show that the following CLIQUE problem is NP-hard:

INPUT: A graph  $G$  and an integer  $k$

OUTPUT: 1 if  $G$  has a clique of size  $k$ , and 0 otherwise

A clique is a collection of mutually adjacent vertices. Use the fact that the following INDEPENDENT SET problem is NP-complete.

INPUT: A graph  $G$  and an integer  $k$

OUTPUT: 1 if  $G$  has an independent set of size  $k$ , and 0 otherwise

An independent set is a collection of mutually nonadjacent vertices.

(b) Show that the following problem is NP-hard:

INPUT: A graph  $G$  **and** an integer  $k$

OUTPUT: 1 if  $G$  has a clique of size  $k$  and an independent set of size  $k$ , and 0 otherwise

Use the fact that both that the CLIQUE and INDEPENDENT SET problems above are NP-complete. This was one of the homework problems.

(c) Show that the following problem is NP-hard:

INPUT: A graph  $G$  and an integer  $k$

OUTPUT: 1 if  $G$  has a clique of size  $k$  **or** an independent set of size  $k$ , and 0 otherwise

Use the fact that both that the CLIQUE and INDEPENDENT SET problems above are NP-complete.