

CS3710 Advanced Topics in AI, Lecture 6

Variable Elimination and Conditioning: Complexity Forecasts.

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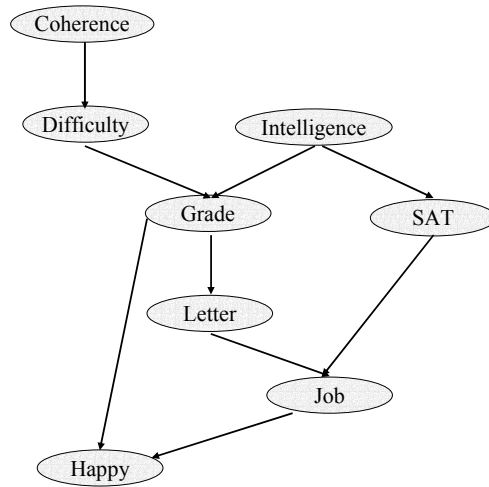
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Outline

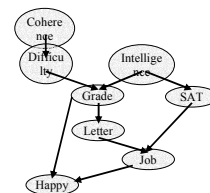
- Recap.
- Factor-Based Elimination.
- Moral Graphs and Triangulation.
- Variable Elimination.
- Induced Graph.
- Correspondence with Tree Decomposition.
- Initial Conclusions.
- Conditioning.
- Final Conclusions.

Recap (Problem).



Recap (Variable Elimination)

$$\begin{aligned}
 p(J) &= \sum_{L,S,G,H,I,D,C} \phi(c)\phi(i)\phi(d,c)\phi(g,i,d)\phi(s,i)\phi(l,g)\phi(j,l,s)\phi(h,g,j) \\
 &= \sum_{L,S,G,H,I,D} \phi(i)\phi(g,i,d)\phi(s,i)\phi(l,g)\phi(j,l,s)\phi(h,g,j) \sum_C \phi(c)\phi(d,c) \\
 &= \sum_{L,S,G,H,I,D} \phi(i)\phi(g,i,d)\phi(s,i)\phi(l,g)\phi(j,l,s)\phi(h,g,j)\tau(c) \\
 &\dots \\
 &= \sum_{S,I} \phi(j,l,s) \sum_G \phi(l,g)\tau(s,g)\tau(g,j) \\
 &= \sum_{S,I} \phi(j,l,s)\tau(l,s,j) \\
 &= \sum_I \tau(l,j) \\
 &= \tau(j)
 \end{aligned}$$



Basic messages

- Variable Elimination is not deterministic.
- The order of elimination governs the overall efficiency.
- Finding an optimal ordering is difficult.
- While different methodologies exist they are all functionally identical.
- The cost of any ordering is exponential in the number of variables that appear in the largest factor.

Factor-Based Elimination.

$$\begin{aligned} p(J) &= \sum_{L,S,G,H,I,D,C} \phi(c)\phi(i)\phi(d,c)\phi(g,i,d)\phi(s,i)\phi(l,g)\phi(j,l,s)\phi(h,g,j) \\ &= \sum_{L,S,G,H,I,D} \phi(i)\phi(g,i,d)\phi(s,i)\phi(l,g)\phi(j,l,s)\phi(h,g,j) \sum_C \phi(c)\phi(d,c) \\ &= \sum_{L,S,G,H,I,D} \phi(i)\phi(g,i,d)\phi(s,i)\phi(l,g)\phi(j,l,s)\phi(h,g,j)\tau(c) \\ &\dots \\ &= \sum_{S,I} \phi(j,l,s) \sum_G \phi(l,g)\tau(s,g)\tau(g,j) \\ &= \sum_{S,I} \phi(j,l,s)\tau(l,s,j) \\ &= \sum_I \tau(l,j) \\ &= \tau(j) \end{aligned}$$

FBE: Trace

Step	Var	Factors Used	New Factor
1	C	$\phi_C(C), \phi_D(D, C)$	$\tau_1(D)$
2	D	$\phi_G(G, I, D), \tau_1(D)$	$\tau_2(G, I)$
3	I	$\phi_I(I), \phi_S(S, I), \tau_2(G, I)$	$\tau_3(G, S)$
4	H	$\phi_H(H, G, J)$	$\tau_4(G, J)$
5	G	$\tau_4(G, J), \tau_3(G, S), \phi_L(L, G)$	$\tau_5(J, L, S)$
6	S	$\tau_5(J, L, S), \phi_J(J, L, S)$	$\tau_6(J, L)$
7	L	$\tau_6(J, L)$	$\tau_7(J)$

FBE: Trace

Step	Var	Factors Used	New Factor	Complexity
1	C	$\phi_C(C)\phi_D(D, C)$	$\tau_1(D)$	2
2	D	$\phi_G(G, I, D)\tau_1(D)$	$\tau_2(G, I)$	3
3	I	$\phi_I(I)\phi_S(S, I)\tau_2(G, I)$	$\tau_3(G, S)$	3
4	H	$\phi_H(H, G, J)$	$\tau_4(G, J)$	3
5	G	$\tau_4(G, J)\tau_3(G, S)\phi_L(L, G)$	$\tau_5(J, L, S)$	4
6	S	$\tau_5(J, L, S)\phi_J(J, L, S)$	$\tau_6(J, L)$	3
7	L	$\tau_6(J, L)$	$\tau_7(J)$	2

Factor-based elimination

- **Ordering Ω :**
 - A permutation of variables for elimination.
- **Factor $\Phi(X, Y) \rightarrow \mathcal{R}$:**
 - A function mapping some set of variables to a real value.
- **$scope[\Phi]$:**
 - The set of variables represented in the factor.
- **$width[\Omega]$:**
 - The scope of the largest factor produced by Ω .

FBE: Steps

- FBE consists of a series of elimination steps.
- Each step is as follows:
 - Select a variable X from the set of variables remaining.
 - Multiply all factors τ where $X \in scope[\tau]$ to produce a new factor ψ .
 - Sum X out of ψ to produce a new factor τ whose scope is ψ minus X .
- Repeat until a single factor $\tau(Y)$ remains where Y is the target variable of our inference.

Complexity

- The complexity of each elimination step is: $O(N_i k_i)$
 - Where:

$$\psi_i = \phi_1 \times \dots \times \phi_{k_i}$$

$$N_i = |scope(\psi_i)|$$

- The complexity of the algorithm for a given ordering Ω is: $O(nN_{\max})$
 - Where:
 - n is the initial number of factors in the graph.
 - $N_{\max} = width[\Omega]$.

FBE: Trace

Step	Var	Factors Used	New Factor
1	C	$\phi_c(C), \phi_D(D, C)$	$\tau_1(D)$
2	D	$\phi_G(G, I, D), \tau_1(D)$	$\tau_2(G, I)$
3	I	$\phi_I(I), \phi_S(S, I), \tau_2(G, I)$	$\tau_3(G, S)$
4	H	$\phi_H(H, G, J)$	$\tau_4(G, J)$
5	G	$\tau_4(G, J), \tau_3(G, S), \phi_L(L, G)$	$\tau_5(J, L, S)$
6	S	$\tau_5(J, L, S), \phi_J(J, L, S)$	$\tau_6(J, L)$
7	L	$\tau_6(J, L)$	$\tau_7(J)$

FBE: Trace.

Step	Var	Factors Used	New Factor	Complexity
1	C	$\phi_C(C)\phi_D(D, C)$	$\tau_1(D)$	2
2	D	$\phi_G(G, I, D)\tau_1(D)$	$\tau_2(G, I)$	3
3	I	$\phi_I(I)\phi_S(S, I)\tau_2(G, I)$	$\tau_3(G, S)$	3
4	H	$\phi_H(H, G, J)$	$\tau_4(G, J)$	3
5	G	$\tau_4(G, J)\tau_3(G, S)\phi_L(L, G)$	$\tau_5(J, L, S)$	4
6	S	$\tau_5(J, L, S)\phi_J(J, L, S)$	$\tau_6(J, L)$	3
7	L	$\tau_6(J, L)$	$\tau_7(J)$	2

Total cost: $width[\Omega] = 4$

$$O(nN_{\max}) = O(8 \times 4) = O(32)$$

Ordering 2

Step	Var	Factors Used	New Factor
1	G	$\phi_G(G, I, D), \phi_L(L, G)\phi_H(H, G, J)$	$\tau_1(I, D, L, J, H)$
2	I	$\phi_I(I), \phi_S(S, I)\tau_1(I, D, L, J, H)$	$\tau_2(D, L, S, J, H)$
3	S	$\phi_J(J, L, S), \tau_2(D, L, S, J, H)$	$\tau_3(D, L, J, H)$
4	L	$\tau_3(D, L, J, H)$	$\tau_4(D, J, H)$
5	H	$\tau_4(D, J, H)$	$\tau_5(D, J)$
6	C	$\tau_5(D, J), \phi_D(D, C)$	$\tau_6(D, J)$
7	D	$\tau_6(D, J)$	$\tau_7(J)$

FBE: Trace.

Step	Var	Factors Used	New Factor	Complexity
1	C	$\phi_G(G, I, D)\phi_L(L, G)\phi_H(H, G, J)$	$\tau_1(I, D, L, J, H)$	6
2	D	$\phi_I(I)\phi_S(S, I)\tau_1(I, D, L, J, H)$	$\tau_2(D, L, S, J, H)$	6
3	I	$\phi_J(J, L, S)\tau_2(D, L, S, J, H)$	$\tau_3(D, L, J, H)$	5
4	H	$\tau_3(D, L, J, H)$	$\tau_4(D, J, H)$	4
5	G	$\tau_4(D, J, H)$	$\tau_5(D, J)$	3
6	S	$\tau_5(D, J)\phi_D(D, C)$	$\tau_6(D, J)$	3
7	L	$\tau_6(D, J)$	$\tau_7(J)$	2

Total cost: $width[\Omega] = 6$

$$O(nN_{\max}) = O(8 \times 6) = O(48)$$

Optimal Permutation.

Optimal Ordering Ω' :

$$\Omega' \text{ s.t. } \forall \Omega': width[\Omega'] < width[\Omega]$$

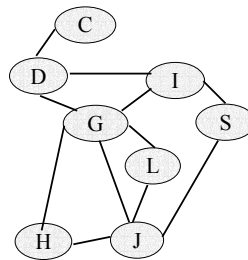
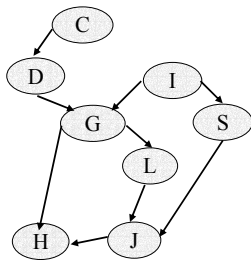
Note:

- The optimal ordering is not guaranteed to be unique.
- Nor is it guaranteed to be less than: $O(nN_{\max})$

Moral Graphs

Moral-graph $H[G]$: of a bayesian network over X is an undirected graph over X that contains an edge between x and y if:

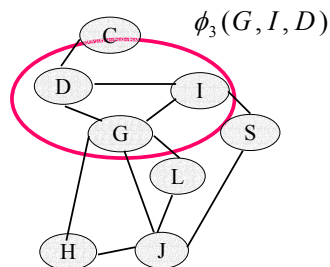
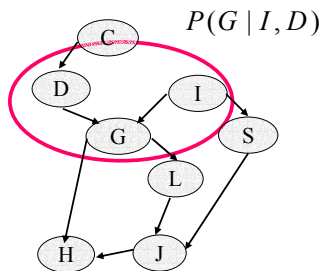
- There exists a directed edge between them in G .
- They are both parents of the same node in G .



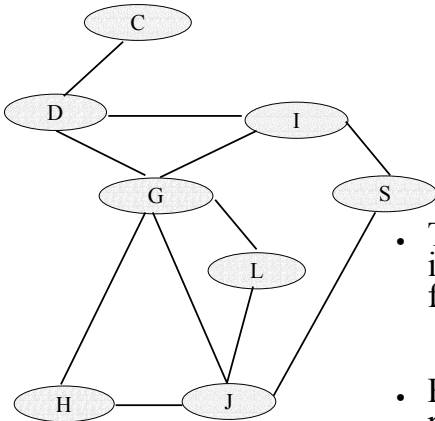
Moral Graphs

Why moralization?

$$\begin{aligned}
 P(C, D, G, I, S, L, J, H) &= \\
 &= P(C)P(D | C)P(G | I, D)P(S | I)P(L | G)P(J | L, S)P(H | G, J) \\
 &= \phi_1(C)\phi_2(D, C)\phi_3(G, I, D)\phi_4(S, I)\phi_5(L, G)\phi_6(J, L, S)\phi_7(H, G, J)
 \end{aligned}$$

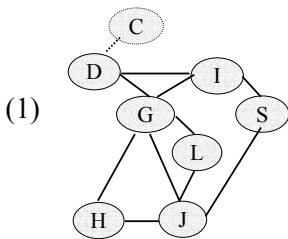


Variable Elimination



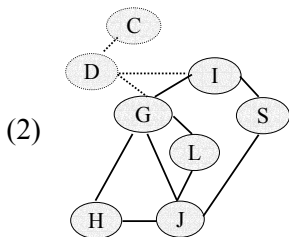
- The variable elimination algorithm is based only on the scope of each factor.
- Factors are distilled from the graph representation.
- Variable Elimination can therefore be viewed as a graph algorithm.

VE: Trace (1)



- Multiply the factors to produce:
 $\phi(D, C) = \phi(D) \times \phi(C)$
- Sum over C to produce:

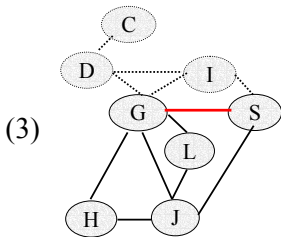
$$\tau(D) = \sum_C \phi(D, C)$$



- Multiply the factors to produce:
 $\phi(D, I, G) = \phi(G, I, D) \times \tau_1(D)$
- Sum over D to produce:

$$\tau_2(G, I) = \sum_D \phi(D, I, G)$$

VE: Trace (2)

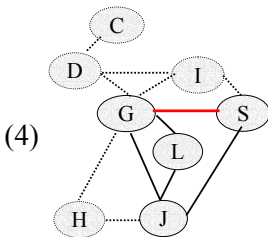


a) Multiply the factors to produce:

$$\phi(I, G, S) = \phi(I) \times \phi(S, I) \times \tau_2(G, I)$$

b) Sum over I to produce:

$$\tau_3(G, S) = \sum_I \phi(I, G, S)$$



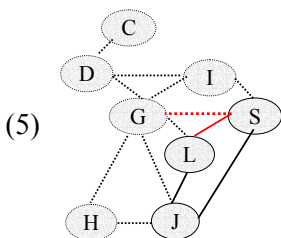
a) Multiply the factors to produce:

$$\phi(H, G, J) = \phi(H, G, J)$$

b) Sum over I to produce:

$$\tau_4(G, J) = \sum_H \phi(H, G, J)$$

VE: Trace (3)

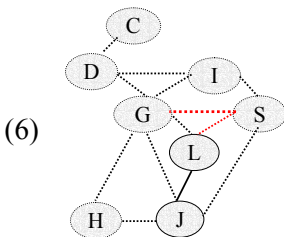


a) Multiply the factors to produce:

$$\phi(G, J, L, S) = \tau_4(G, J) \times \tau_3(G, S) \times \phi(L, G)$$

b) Sum over I to produce:

$$\tau_5(J, L, S) = \sum_G \phi(G, J, L, S)$$



a) Multiply the factors to produce:

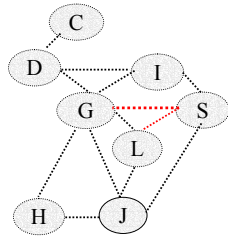
$$\phi(J, L, S) = \tau_5(J, L, S) \times \phi(J, L, S)$$

b) Sum over I to produce:

$$\tau_6(J, L) = \sum_S \phi(J, L, S)$$

VE: Trace (4)

(5)



a) Multiply the factors to produce:

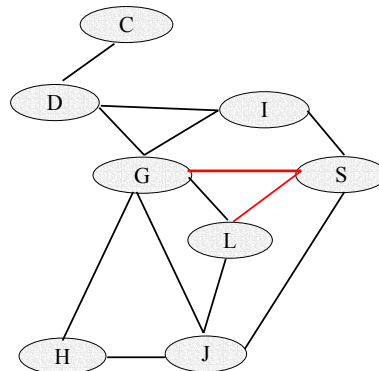
$$\phi(J, L) = \tau_6(J, L)$$

b) Sum over I to produce:

$$\tau_7(J) = \sum_L \phi(J, L)$$

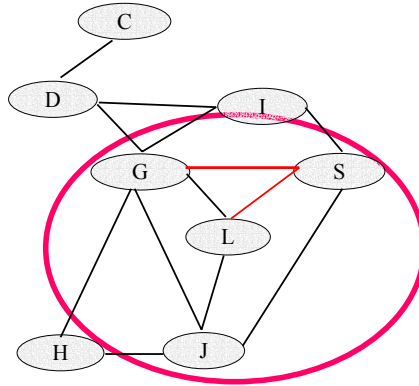
Variable elimination: Induced Graph

- Induced Graph G' : An undirected graph over X where y and z are connected if they both appear in some intermediate elimination factor of Ω .
- Every factor generated during Ω appears as a subclique of the graph.



Variable elimination: Induced Graph

- Induced Graph G' : An undirected graph over X where y and z are connected if they both appear in some intermediate elimination factor of Ω .
- Every factor generated during Ω appears as a subclique of the graph.
- The size of the largest clique governs the computation.**



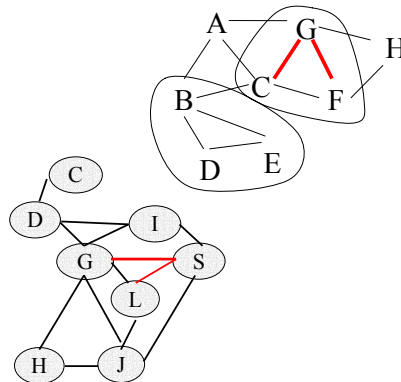
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2	D	$\phi_G(G, I, D), \tau_1(D)$	$\tau_2(G, I)$
3	I	$\phi_I(I), \phi_S(S, I), \tau_2(G, I)$	$\tau_3(G, S)$
4	H	$\phi_H(H, G, J)$	$\tau_4(G, J)$
5	G	$\tau_4(G, J), \tau_3(G, S), \phi_L(L, G)$	$\tau_5(J, L, S)$
6	S	$\tau_5(J, L, S), \phi_J(J, L, S)$	$\tau_6(J, L)$
7	L	$\tau_6(J, L)$	$\tau_7(J)$

lence

a1	b1	0.5
a1	b2	0.2
a2	b1	0.1
a2	b2	0.3
a3	b1	0.2
a3	b2	0.4

b1	c1	0.1
b1	c2	0.6
b2	c1	0.3
b2	c2	0.4

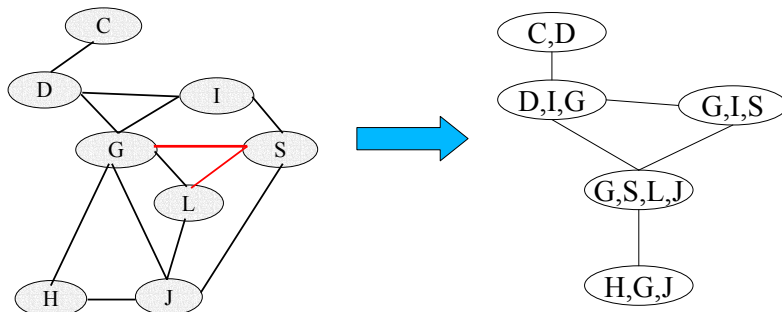
a1	b1	c1	0.5*0.1
a1	b1	c2	0.5*0.6
a1	b2	c1	0.2*0.3
a1	b2	c2	0.2*0.4
a2	b1	c1	0.1*0.1
a2	b1	c2	0.1*0.6
a2	b2	c1	0.3*0.3
a2	b2	c2	0.3*0.4
a3	b1	c1	0.2*0.1
a3	b1	c2	0.2*0.6
a3	b2	c1	0.4*0.3
a3	b2	c2	0.4*0.4



- $Tree-Width = best\ Max\ Factor-Width - 1 = best\ Max\ Clique-Size - 1$
- the tree-width defined by transformation of the undirected graph to the best factor tree is the same

Induced Graph to Tree Decomposition

- For each Clique in the induced graph:
 - Collapse the clique into a compound factor node containing the member variables.
 - Draw a link between each pair of nodes that share a member variable.

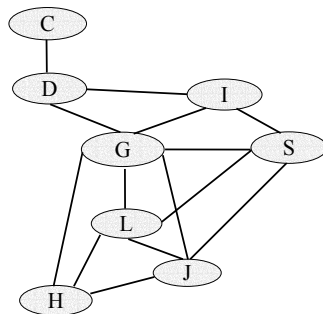


The Bad News

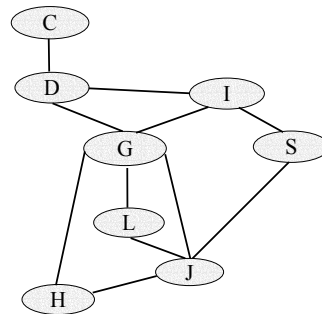
- Deciding whether or not there exists an ordering Ω s.t. $Width(\Omega) < N$ is *NP-Hard*.
 - *No matter what method is used.*
- At best the inference is still exponential in the treewidth of the factor graph

Chordal graphs and Triangulations

Chordal Graph: an undirected graph G whose minimum cycle contains 3 vertices.



Chordal.



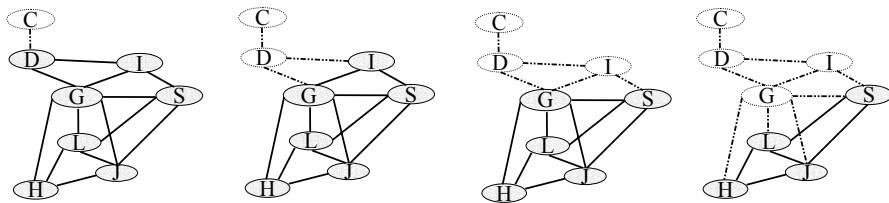
Not Chordal.

Triangulation

- The process of converting a graph G into a chordal graph is called Triangulation.
- The induced graph is:
 - 1) Guaranteed to be chordal.
 - 2) Not guaranteed to be optimal.
- There exist exact algorithms for minimal chordal graphs, and heuristic methods with a guaranteed upper bound.

Chordal Graphs: Good News

- If the moralized graph of our original network is chordal:
 - There exists an elimination ordering that adds no edges.
 - The minimal induced width of the graph is the size of the largest clique - 1.



Chordal Graphs: Bad News

- Given a minimum triangulation for a graph G , we can carry out the variable-elimination algorithm in the minimum possible time.
- However, finding the minimal triangulation is NP-Hard.
 - Time is exponential in terms of the largest clique (factor) in G .

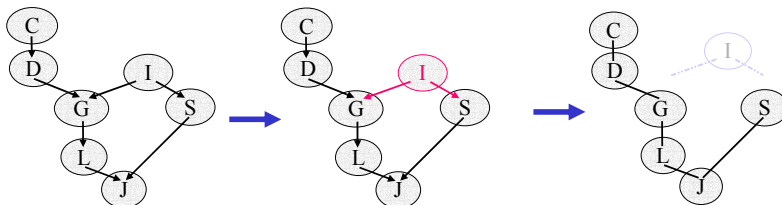
Initial Conclusions

- We cannot escape **exponential costs in the treewidth.**
- But in many graphs the treewidth is much smaller than the total number of variables
- Still a problem: Finding the optimal decomposition is hard
 - But, paying the cost up front may be worth it.
 - Triangulate once, query many times.
 - Real cost savings if not a bounded one.

Conditioning

- Up until now we have ignored some independences
- Assume the Student network from Koller and Friedman

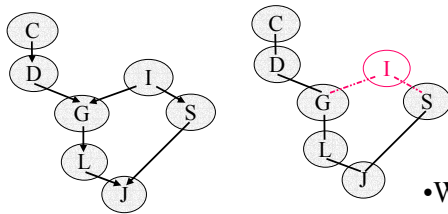
$$\begin{aligned}
 P(J) &= \sum_{c,d,i,g,s,l} P(C = c, D = d, I = i, G = g, S = s, L = l, J) \\
 &= \sum_{c,d,i,g,s,l} P(I = i) P(C = c, D = d, G = g, S = s, L = l, J | I = i) \\
 &= \sum_i P(I = i) \sum_{c,d,g,s,l} P(C = c, D = d, G = g, S = s, L = l, J | I = i) \\
 &= \sum_i P(I = i) P(J | I = i)
 \end{aligned}$$



Conditioning

Concept:

- Fix some variable I s.t. $I=i$ and remove it from the graph.
- Recalculate neighboring probabilities.
- Execute VE with the reduced graph.

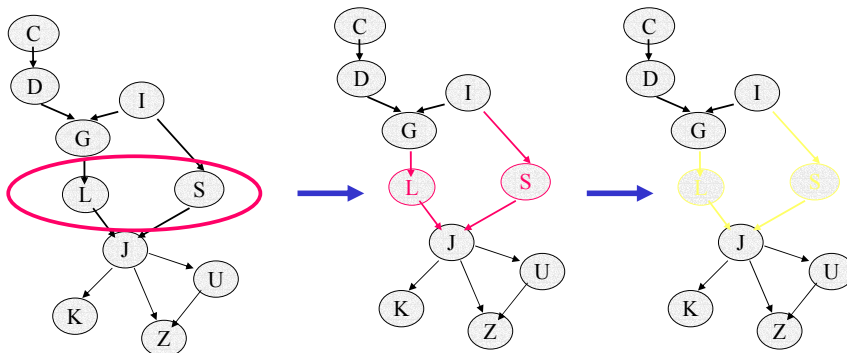


- We need to do a separate inference for every value i of I
- But the structure of the network is much simpler

Conditioning

General conditioning works for sets of variables.

Assume:



Conditioning conclusions

- Conditioning on some variables may simplify the structure of the remaining network
 - The network may break into small pieces.
- Variable elimination may be easier to perform on the new network